

2.5.2 Offset Binary representation (Excess-K) • Offset Binary is where one subtracts K (usually half the largest possible number) from the representation to get the value.

- Has the advantage that the number sequence from the most negative to the most positive is a simple binary progression, which makes it a natural for binary counters.
- note that the MSB still carries the sign information.
- Excess K is used in conjunction with floating point representations for the exponent. We will meet this again shortly. • A note on arithmetic in Excess K:

Assume a, b, c are three values:

$$\begin{aligned} (a + b) &= c && \text{(values)} \\ (a + k) + (b + k) &&& \text{(representations)} \\ &= (a + b) + 2k \\ &= (c + k) + k \end{aligned}$$

Rewriting A, B, C in Excess-K representations:

$$\begin{aligned} A + B &= C + k \\ C &= (A + B) - k \end{aligned}$$

- Try this for  $(-1) + (+1) = 0$

### 2.5.3 2's complement

- 2's complement represents the method most widely used for integer computation.
- Positive numbers are represented in simple unsigned binary.
- The system is rigged so that a negative number is represented as the binary number that when added to a positive number of the same magnitude gives zero.
- To get the two's complement, first take the ones complement, then add one.

### 2.5.4 1's complement

- Exchange all the 1's for 0's and vice versa.

value	1's complement	2's complement
+7	0111	0111
+6	0110	0110
+5	0101	0101
+4	0100	0100
+3	0011	0011
+2	0010	0010
+1	0001	0001
0	0000	0000
-1	1110	1111
-2	1101	1110
-3	1100	1101
-4	1011	1100
-5	1010	1011
-6	1001	1010
-7	1000	1001
-8	-	1000
-0	1111	-

Binary Value	1's-Complement	2's-complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	-7	-8
1001	-6	-7
1010	-5	-6
1011	-4	-5
1100	-3	-4
1101	-2	-3
1110	-1	-2
1111	-0	-1

## 2.6 Performing Arithmetic

### 2.6.1 In 1's complement

Some examples of arithmetic with 1's complement.

Example 2.6.1 (addition)

$$\begin{array}{r} 0011 \quad (+3) \\ +0010 \quad (+2) \\ \hline 0101 \quad (+5) \end{array}$$

Example 2.6.2 (subtraction)

$$\begin{array}{r} 0011 \quad (+3) \\ +1101 \quad (-2) \\ \hline (1)0000 \quad (0?) \end{array}$$

Example 2.6.3 (subtraction from a negative)

$$\begin{array}{r} 1100 \quad (-3) \\ +1101 \quad (-2) \\ \hline (1)1001 \quad (-6?) \end{array}$$

The solution is to wrap the carry back in to the LSB.

Exercise 2.6.4 Can you explain why this works?

### 2.6.2 In 2's complement

The Arithmetic operations are perhaps easiest in 2's complement.

- To add ... just like in any other base.

Example 2.6.5 (addition 5 + (-2):)

$$\begin{array}{r}
 0101 \quad (+5) \\
 +1110 \quad (-2) \\
 \hline
 0011 \quad (+3)
 \end{array}$$

- To subtract B from A take the 2's complement of B and add to A.

Example 2.6.6 (subtraction 2 - 5:)

$$\begin{array}{r}
 0010 \quad (+2) \quad (2 + (-5)) \\
 +1011 \quad (-5) \text{ since } +5 = 0101: \\
 \hline
 1101 \quad (-3)
 \end{array}$$

value	Sign Magnitude	Offset Binary	2's complement
+7	0111	1111	0111
+6	0110	1110	0110
+5	0101	1101	0101
+4	0100	1100	0100
+3	0011	1011	0011
+2	0010	1010	0010
+1	0001	1001	0001
0	0000	1000	0000
-1	1001	0111	1111
-2	1010	0110	1110
-3	1011	0101	1101
-4	1100	0100	1100
-5	1101	0011	1011
-6	1110	0010	1010
-7	1111	0001	1001
-8	-	0000	1000
-0	1000	-	-

Binary Value	Sign Magnitude	Offset Binary	2's complement
0000	0	-8	0
0001	1	-7	1
0010	2	-6	2
0011	3	-5	3
0100	4	-4	4
0101	5	-3	5
0110	6	-2	6
0111	7	-1	7
1000	-0	0	-8
1001	-1	1	-7
1010	-2	2	-6
1011	-3	3	-5
1100	-4	4	-4
1101	-5	5	-3
1110	-6	6	-2
1111	-7	7	-1

- Multiplication also works right in 2's complement. Long multiplication reduces to shifts and adds
- We have implicitly used the concept of carry. In particular we dropped/ignored the carry bit in the case of the two's complement number representation.

Example 2.6.7 ( $3 - 3 =$ )

$$\begin{array}{r}
 0011 \quad (3) \\
 +1101 \quad (-3) \\
 \hline
 (1)0000 \quad (0)
 \end{array}$$

(c.f. above 1's complement example)

- It should be clear that for the unsigned binary the carry has relevance.

Example 2.6.8 (3 + 13 =)

$$\begin{array}{r} 0011 \quad (3) \\ +1101 \quad (+13) \\ \hline (1) 0000 \quad (16) \end{array}$$

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- We can also perform subtraction directly and still ignore the carry/borrow bit.

Example 2.6.9 (2's complement borrow)

$$\begin{array}{r} 0011 \quad (3) \\ -0100 \quad (-4) \\ \hline (1) 1111 \quad (-1) \end{array}$$

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- However, for subtraction with the unsigned binary the borrow is important, particularly in multiple word operations.

Example 2.6.10 (multiple word addition/carry)

$$\begin{array}{r} 0011 \quad 0011 \quad (51) \\ +0000^11101 \quad (13) \\ \hline 0100 \quad 0000 \quad (64) \end{array}$$

Example 2.6.11 (multiple subtraction/borrow)

$$\begin{array}{r} 0100 \quad 0000 \quad (64) \\ -0000^11101 \quad (-13) \\ \hline 0011 \quad 0011 \quad (51) \end{array}$$

- These still could represent two's complement. but low order words must treated as if unsigned.

